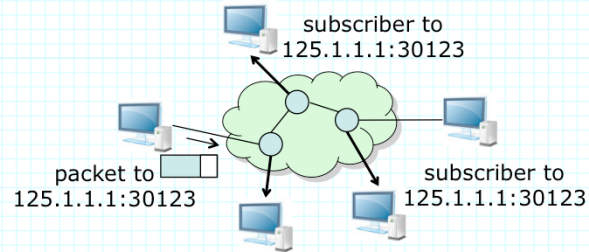


19. Multicast Communication

- Multicast within a subnet
- Reverse path forwarding
- IGMP
- PIM
- Scalable multicast forwarding

Jon Turner

Multicast in a Subnet



- Within switched subnet, multicasting relies on Ethernet
 - » IP multicast address embedded in Ethernet multicast address
 - more precisely, low 23 bits are embedded
 - » switches deliver multicast packets to all hosts (simple case)
- Receiving application configures multicast socket (UDP)
 - » bind to specific port, join multicast group
- Sending app sends UDP packet on multicast socket
 - » use multicast address/port, set TTL to 1

Basic Multicast Receiver

```
/** usage: Rcvr mcAdr port iface */
import java.io.*; import java.net.*;
public class Rcvr {
    public static void main(String args[]) throws Exception {
        // form multicast group from command-line arguments
        int port = Integer.parseInt(args[1]);
        InetSocketAddress group = new InetSocketAddress(
            InetAddress.getByName(args[0]),port);
        // open multicast socket and join group
        MulticastSocket sock = new MulticastSocket(port);
        NetworkInterface iface = NetworkInterface.getBy_name(args[2]);
        sock.joinGroup(group,iface);
        // create buffer and packet
        byte[] inBuf = new byte[1000];
        DatagramPacket inPkt = new DatagramPacket(inBuf,inBuf.length);
        while (true) {
            // receive packet and print payload
            sock.receive(inPkt);
            System.out.println(new
                String(inBuf,0,inPkt.getLength(),"US-ASCII"));
        }
    }
}
```

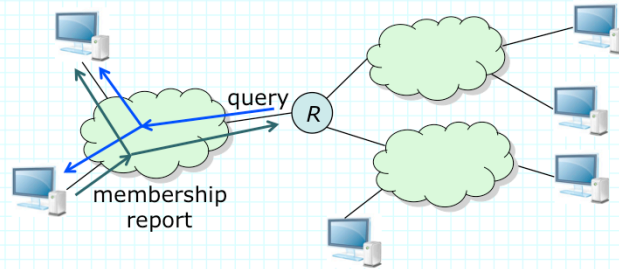
Basic Multicast Sender

```
/** usage: Sender mcAdr port iface ttl message */
import java.io.*; import java.net.*;
public class Sender {
    public static void main(String args[]) throws Exception {
        // from multicast group address from command-line arguments
        int port = Integer.parseInt(args[1]);
        InetSocketAddress group = new InetSocketAddress(
            InetAddress.getByName(args[0]),port);

        // open multicast socket, set interface and time to live
        MulticastSocket sock = new MulticastSocket();
        sock.setNetworkInterface(NetworkInterface.getByName(args[2]));
        sock.setTimeToLive(Integer.parseInt(args[3]));

        // create buffer and packet to send
        byte[] outBuf = args[4].getBytes("US-ASCII");
        DatagramPacket outPkt = new DatagramPacket(
            outBuf,outBuf.length,group);
        sock.send(outPkt); // send packet to group
        sock.close();
    }
}
```

Internet Group Multicast (RFC 2236)



- To forward among subnets, router needs to know about multicast subscribers
 - » router multicasts *membership query* packets
 - » hosts respond with *membership reports* after random delay
 - suppress duplicate reports; usually just one actual report per group
- Router forwards multicast packets among networks based on membership information
 - » subscriptions timeout if not renewed periodically

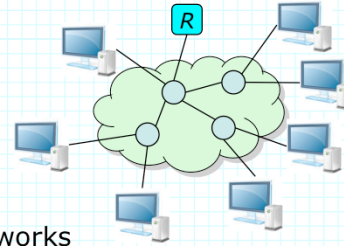
IGMP Details

- Message fields
 - » type – query, report, leave group
 - » max response time – max wait time before responding to query
 - » group address – subject multicast addr; 0 for “general query”
- On each subnet, one router plays role of “Querier”
 - » router with smallest IP address acts as querier
 - » issues periodic general membership queries
 - sent on all systems multicast address: 224.0.0.1
- Hosts send reports to subject multicast address
 - » in response to general queries or when first joining a group
- Leave group message sent by host when leaving
 - » sent to all routers group (224.0.0.2)
 - » querier responds with a group-specific query
 - verify that no other group members still out there

IGMP Snooping (RFC 4541)

■ Objective

- » use group membership to limit propagation of multicast packets
- » learn which hosts are in which group by snooping on IGMP
- » invented by switch vendors to limit multicast traffic in large switched networks



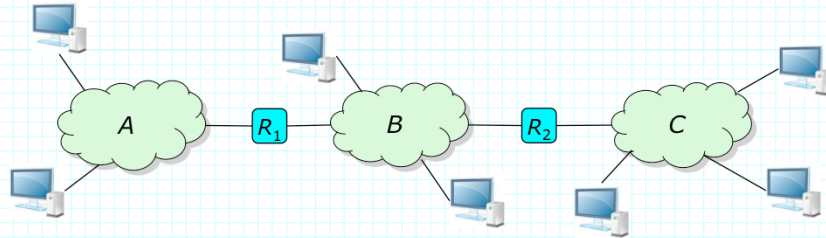
■ Suppression of duplicate membership reports by IGMP interferes with snooping

- » so send report packets only to routers, not to hosts
 - implies switches must be able to detect which ports go to routers

■ Forwarding based on group membership

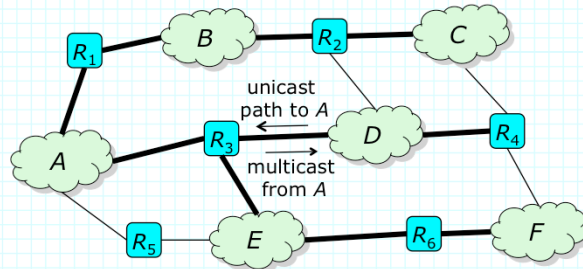
- » forwarding can use either Ethernet or IP addresses
 - IP addresses preferred due to non-exact multicast address mapping
 - but not an option for simple Ethernet switches

Multicast Issues in Larger Networks



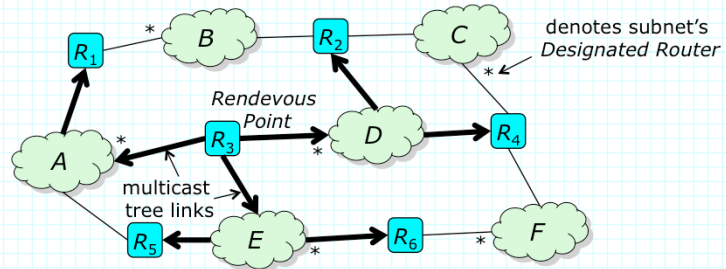
- IGMP is not always enough
 - » suppose a multicast has members in subnets A and C, but not B
 - routers will not forward packets across B
 - » can work if group members are in contiguous subnets
- Multiple paths among subnets can produce multicast duplicates
 - » need to restrict routing of each multicast group to a tree
 - » requires routers to maintain forwarding information for each multicast group

Reverse Path Forwarding



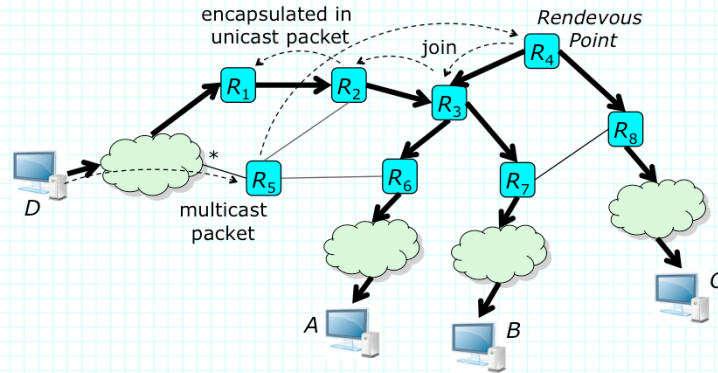
- Efficient distribution of multicast packets requires that packets be forwarded along a tree
 - » in principle, any tree joining subnets can be used and each multicast can use a different tree
 - » internet multicast generally uses reverse of unicast routes
 - so routing based on combination of (source adr, multicast adr)
 - means multicast with many senders requires more routing state
 - more efficient to use *shared tree* through “rendevouz point”

Multicast Using Shared Tree (PIM)



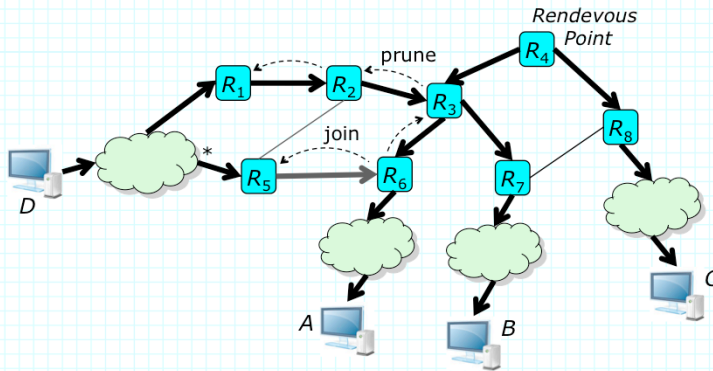
- Rendezvous point (RP) serves as multicast tree root
 - » different groups have different RPs, routers can determine assigned RP for each multicast address
- When *Designated Router* (DR) discovers new subscriber
 - » sends join message to RP, creating multicast state along path
- Sending hosts just send packets to multicast address
 - » initially, DR may forward to RP inside unicast packet

Sending to Multicast Group



- Non-subscribing sender's packets sent to RP by DR
 - » RP extracts multicast packet and forwards on shared tree
- RP may initiate join back towards source
 - » adds multicast routing state along path, merges with RP tree
 - » when native packets reach RP, it tells DR to stop encapsulating

Switch to More Direct Path



- Router in tree may choose to switch to more direct path
 - » sends join back towards source
 - » after packets arrive on direct path, prune links in original path from source
- Note: expensive for multicasts with many senders

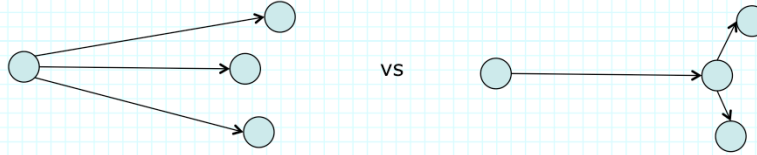
More About PIM (RFC 4601)

- PIM stands for Protocol Independent Multicast
 - » PIM can use routing data from any routing protocol
 - e.g. unicast routing state, or multicast BGP
- Routers must be able to determine RP for any address
 - » can use statically configured mapping based on address ranges
 - » or, a PIM domain can use a *Bootstrap Router* (BSR) to distribute information used to map a multicast address to an RP
 - domain has a defined set of RPs distributed by BSR
 - PIM routers use a hash function to select an RP in this set
 - e.g. if N candidate RPs, $\text{hash}(\text{multicast address}) \bmod N$ selects one
- PIM generally used only within single routing domain
 - » RFC omits essential details for inter-domain routing
 - e.g. how are RPs coordinated among different domains
 - » BGP extensions may eventually enable inter-domain multicast

IGMPv3 (RFC 3376)

- Adds support for source-specific multicast using “filters”
 - » include-filter specifies a list of acceptable senders
 - multicast packets are delivered to socket only if sender is on list
 - » exclude-filter specifies a list of blocked senders
 - » hosts have “interface filters” derived from socket filters
 - » not supported by java multicast sockets; linux does support it
- Queries and membership reports are extended to handle “filtered multicasts”
 - » new source-specific query that includes a sender list for target destination address
 - » reports can cover multiple groups and includes a filter per group
- IGMPv3 ignores membership reports from other hosts
 - » so all IGMPv3 hosts in a multicast group will respond to a query
- Can operate in “compatibility mode” with older versions

Which Trees are Best?



- Reverse path forwarding routes multicast packets on shortest paths
 - » tends to copy packets near the source
 - » copying near destination reduces # of “long-distance packets”
- Best multicast route is a *Steiner tree*
 - » Steiner tree for a vertex subset S is shortest subtree that includes all vertices in S
 - » problem is NP-complete, so no efficient algorithm known for finding optimal solution
 - » good approximations based on minimum spanning tree

Other Issues with IP Multicast

- Separate protocol needed to reserve network capacity
 - » RSVP protocol (RFC 2205) can be used for this purpose
 - » but cannot select route based on required network capacity
 - QoS routing has been discussed but never implemented
- No “private” multicast groups
 - » no way to get short-term exclusive use of a multicast address
 - » no way to limit subscribers or limit who can send to an address
 - makes it awkward to use for multi-party teleconferencing
 - privacy requires encryption
- Largely unavailable in public internet
 - » unresolved technical issues for inter-domain routing
 - » no economic benefit to network providers
 - multicast senders benefit as they need not duplicate traffic
 - but ISPs lose money if senders reduce internet access link rate